

ElastO: Efficient Maintenance of Scalable Overlays for Topic-based Pub/Sub under Churn

[Submission to Middleware '14]

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Joint work with Roman Vitenberg and Hans-Arno Jacobsen

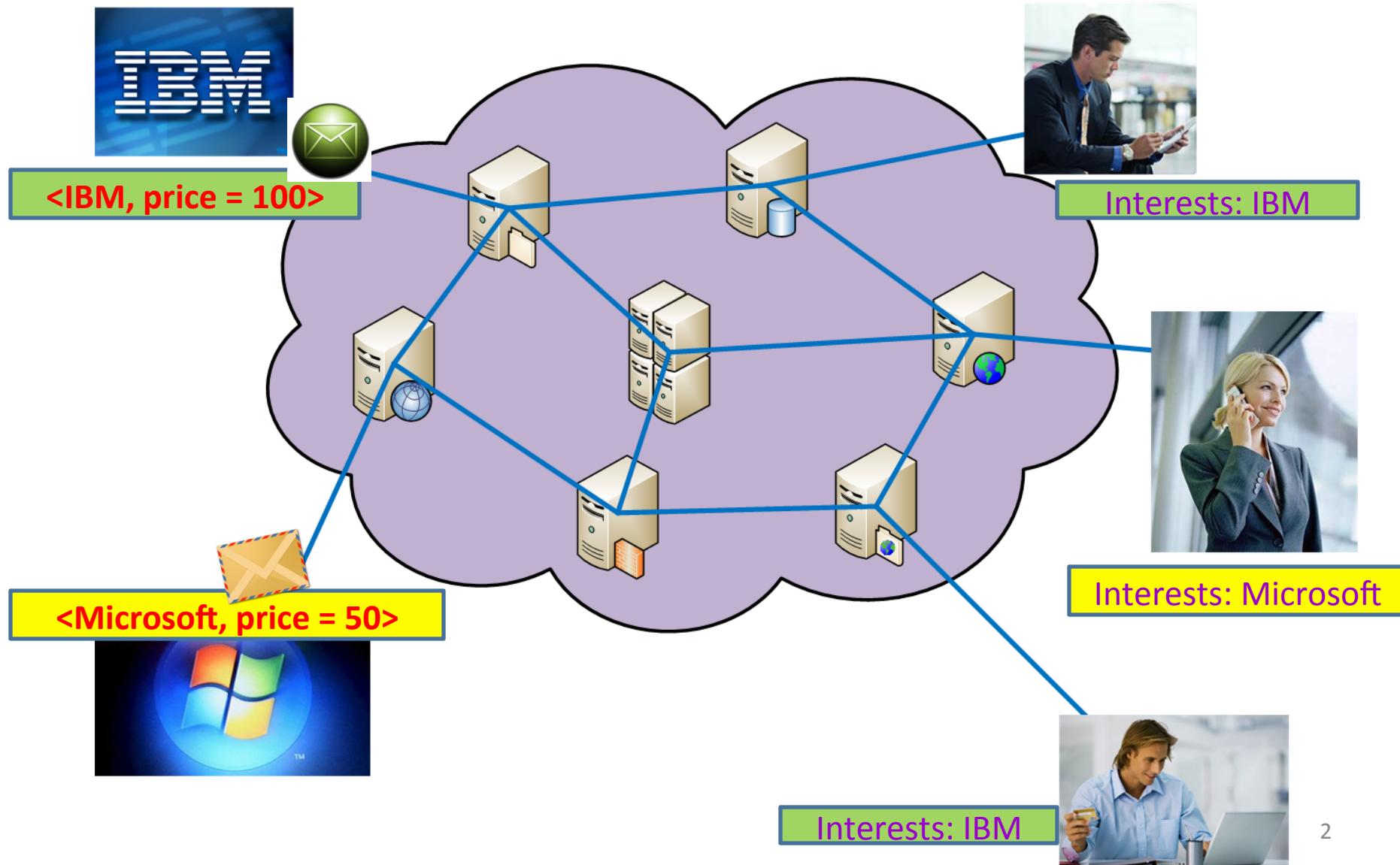


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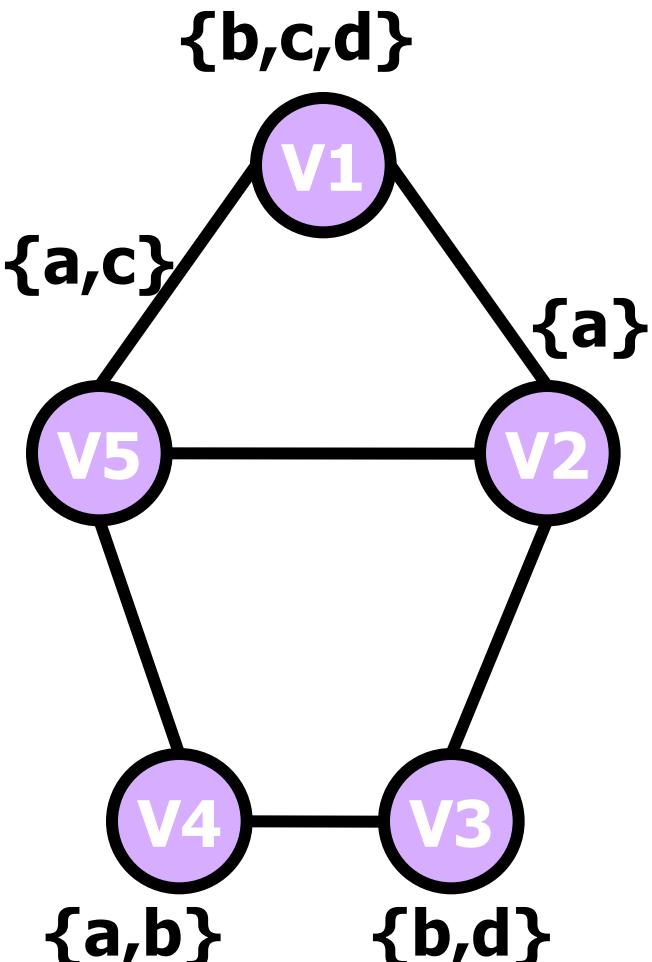


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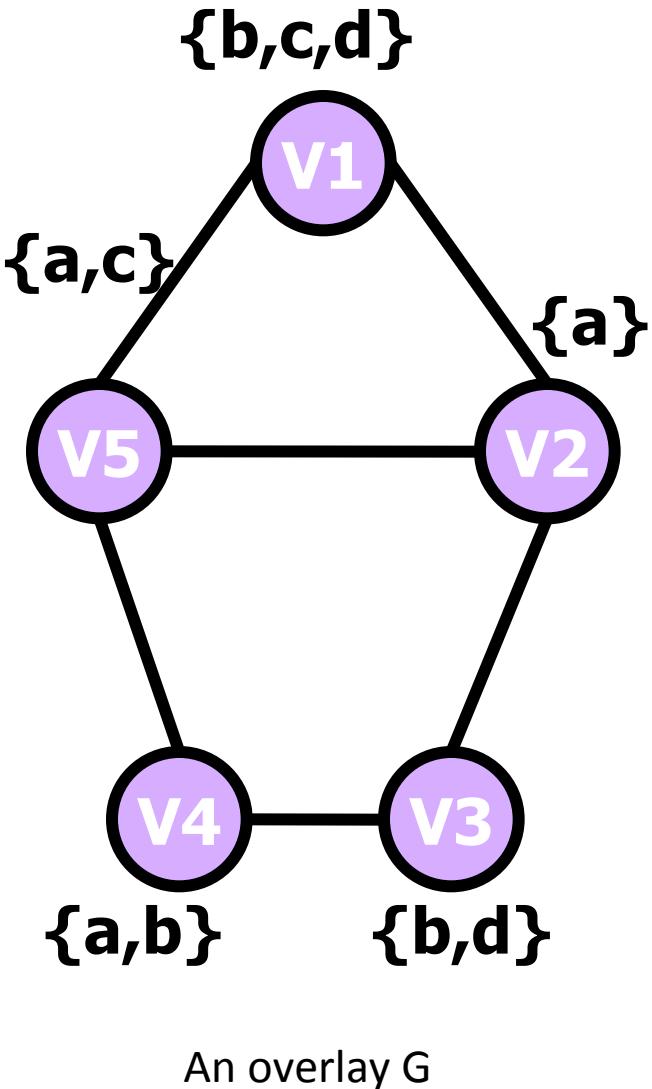


Topic-connected overlay (TCO)



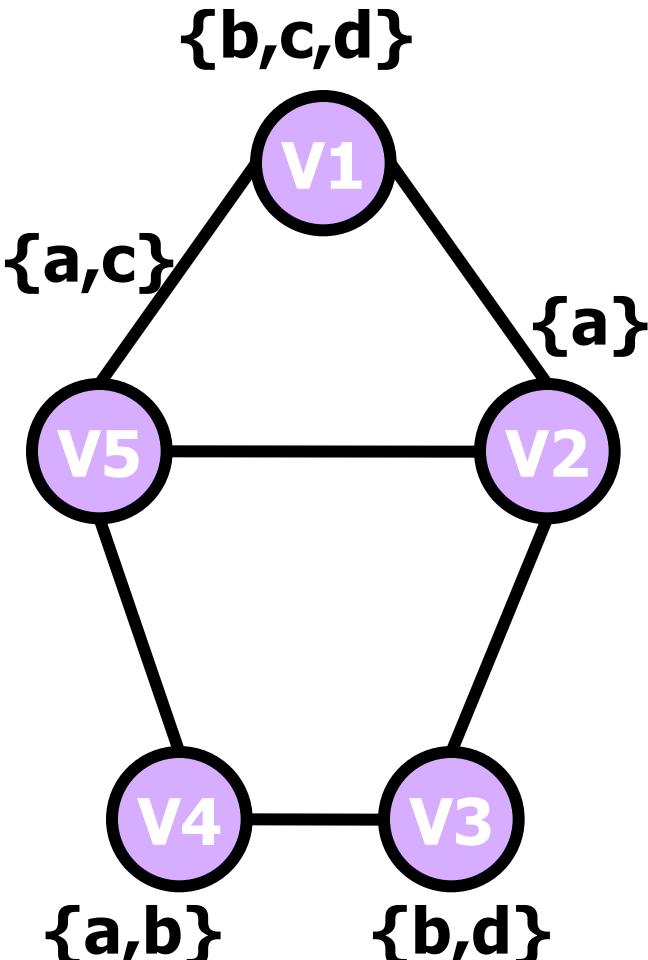
An overlay G

Topic-connected overlay (TCO)



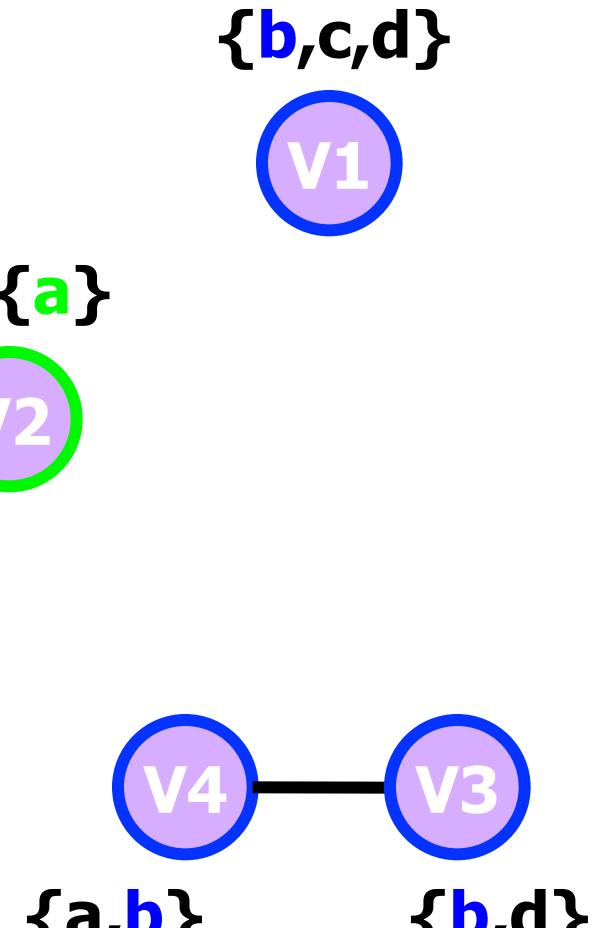
G_a is topic-connected
with one *TC-component*

Topic-connected overlay (TCO)



An overlay G

G_a is topic-connected
with one *TC-component*



G_b is NOT topic-connected
with two *TC-components*

Approaches to build pub/sub TCO

		Knowledge	Churn Handling	Time	Average degree	Maximum degree
Centralized algorithms	Low-ODA	Global	X	Slow	$O(\rho \ln V T)$	$O(V /\rho \ln V T)$
	GM				$O(\ln V T)$	$\Theta(V)$
	MinMax-ODA				$\Theta(V)$	$O(\ln V T)$
	DC				$O(p \ln V T)$	$\Theta(V)$
	DCBR-M				$\Theta(V)$	$O(\eta + \ln V T)$
	GM2				$O(U + \ln V T)$	$\Theta(V)$
Decentralized protocols	SpiderCast, PolderCast, StAN, etc.	Local/Global	✓	Fast	Unknown	Unknown

Our hybrid solution: ElastO

		Knowledge	Churn Handling	Time	Average degree	Maximum degree
Hybrid	ElastO	Local	✓	Fast	$\approx O(\rho \ln V T)$	$\approx O(V /\rho \ln V T)$
Centralized algorithms	Low-ODA	Global	✗	Slow	$O(\rho \ln V T)$	$O(V /\rho \ln V T)$
	GM				$O(\ln V T)$	$\Theta(V)$
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Requirements

A. Overlay quality

- TCO, low node degrees, small diameters
- Close to centralized algorithms

B. Responsiveness

- Comparable to decentralized protocols

C. Scalability of the decentralized solution

- No centralized control
- Partial and local view
- Churn handling only impacts a small portion of nodes

D. Fairness and load balancing

- Computation, communication, storage

E. Reliability against concurrent churn events

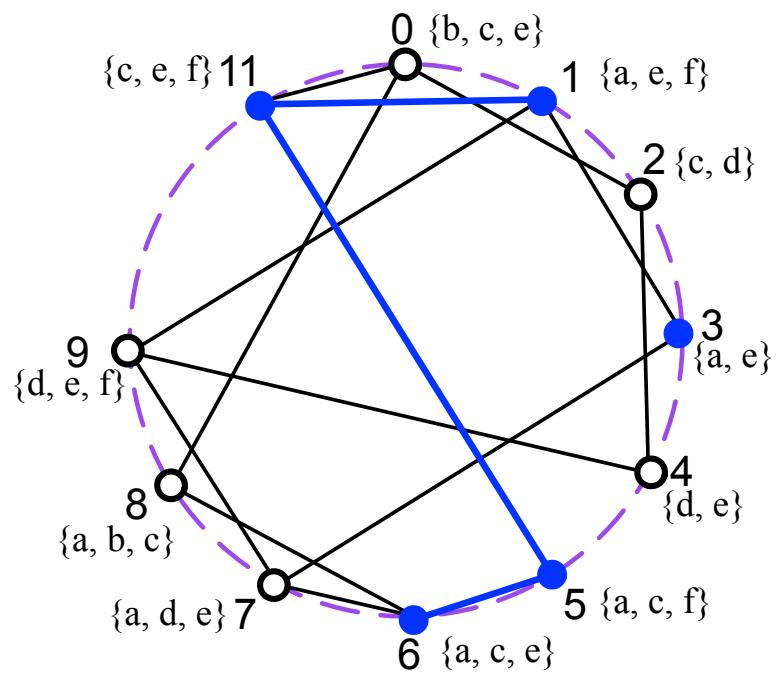
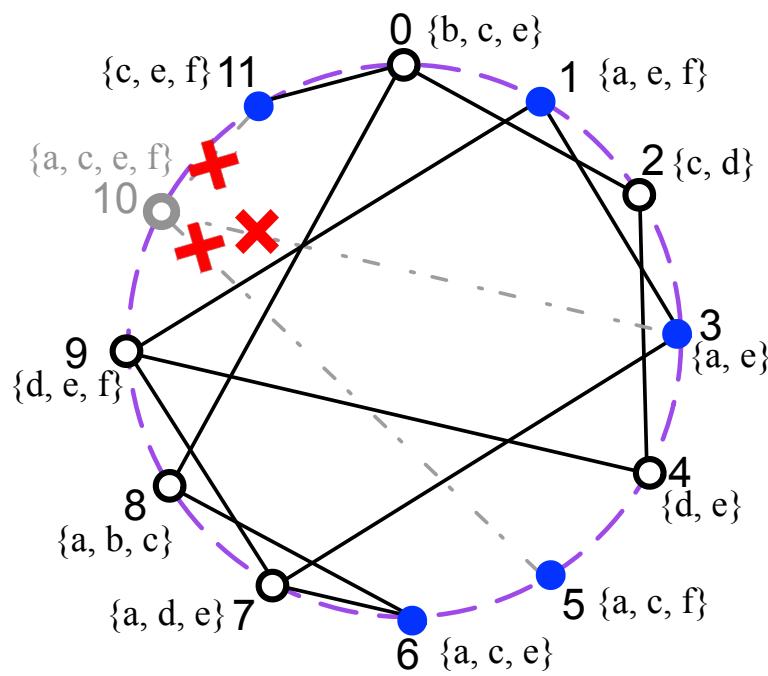
Our contributions - ElastO

- Complete architecture and protocol design
- **Shadow**-based strategies
 - TCO recovery under churn
- **Primary-backup** mechanism
 - support shadow sets in decentralized systems
- Gossip-based peer sampling service with unique **local view selection** algorithms
 - build and deploy backup sets for all nodes
- Comprehensive evaluation

Shadow-based strategies

Upon each churn event

- Select a *proper* subset of nodes, namely the **shadow set**, for TCO recovery



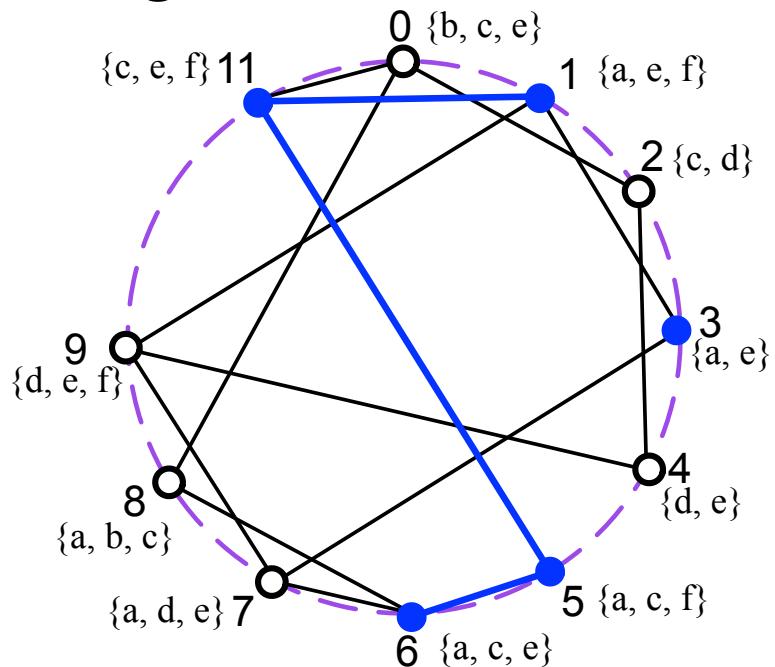
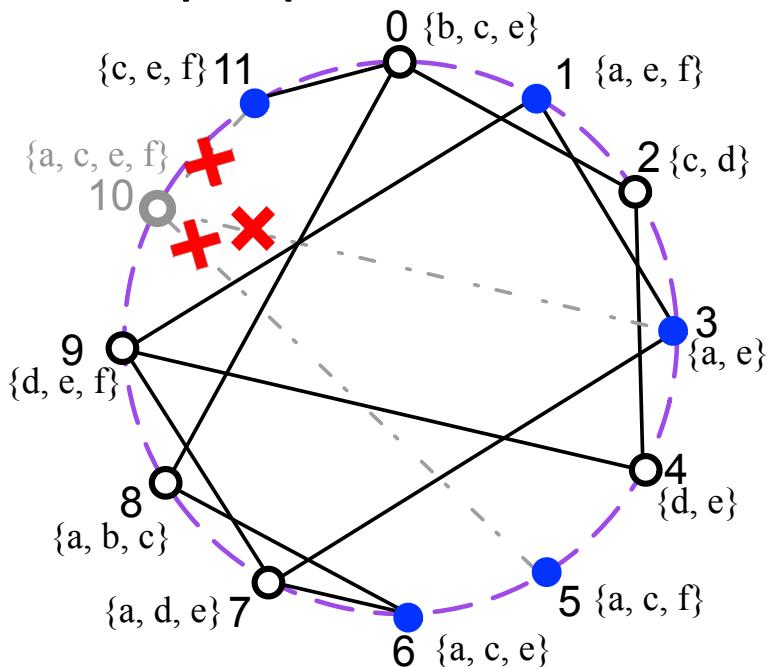
Pre-compute the shadow set primary-backup mechanism

- ElastO needs to remember

1) TCO neighbors, $v_{10}.N = \{v_3, v_5, v_{11}\}$

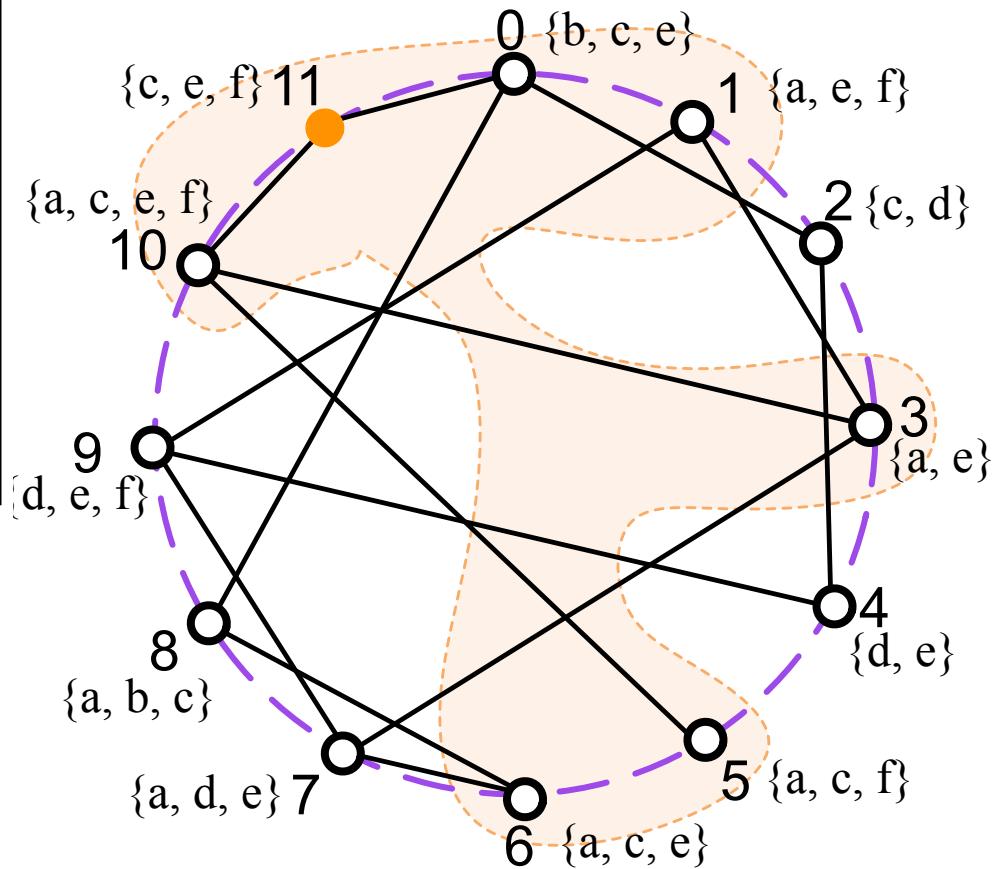
2) backup set, $B(v_{10}) = \{v_1, v_6\}$

in preparation of v_{10} 's leaving



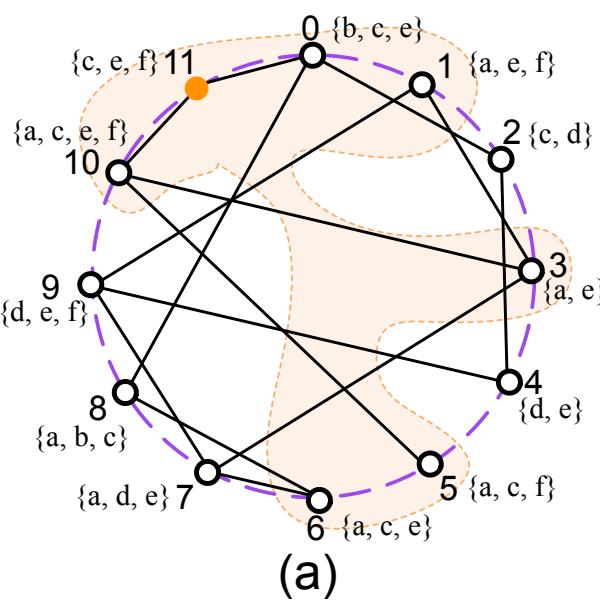
Local view of each node

<i>.self</i>	v_{11}
\mathcal{N}	$\{v_0, v_{10}\}$
\mathcal{D}	\emptyset
$.succ[1]$	v_0
$.pred[1]$	v_{10}
$.pneigh[1]$	$\{v_3, v_5, v_{11}\}$
$.backup[1]$	$\{v_1, v_6\}$



Example: initially stable

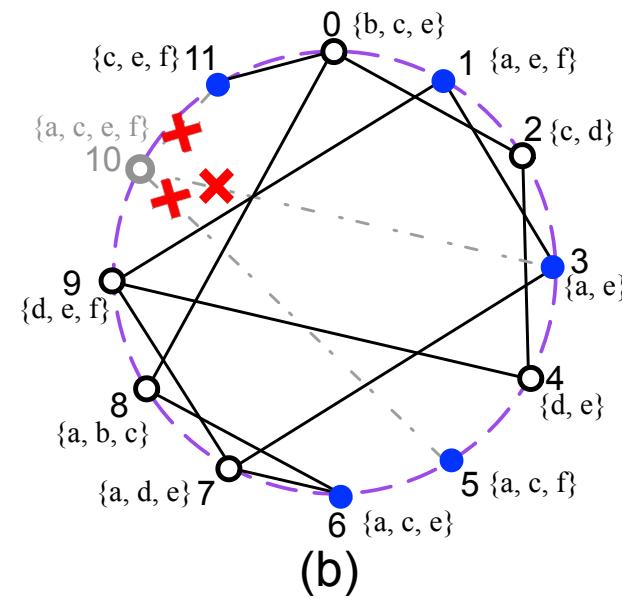
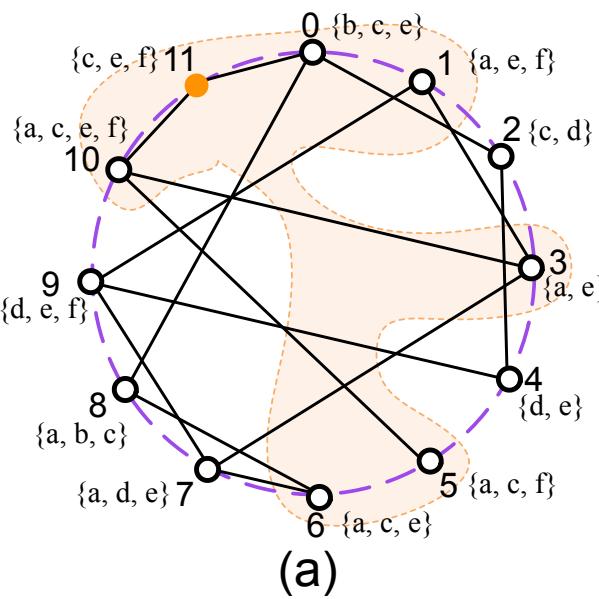
.self	v_{11}
. \mathcal{N}	$\{v_0, v_{10}\}$
. \mathcal{D}	\emptyset
.succ[1]	v_0
.pred[1]	v_{10}
.pneighb[1]	$\{v_3, v_5, v_{11}\}$
.backup[1]	$\{v_1, v_6\}$



Example: node departure

<i>.self</i>	v_{11}
\mathcal{N}	$\{v_0, v_{10}\}$
\mathcal{D}	\emptyset
$.succ[1]$	v_0
$.pred[1]$	v_{10}
$.pneigh[1]$	$\{v_3, v_5, v_{11}\}$
$.backup[1]$	$\{v_1, v_6\}$

<i>.self</i>	v_{11}
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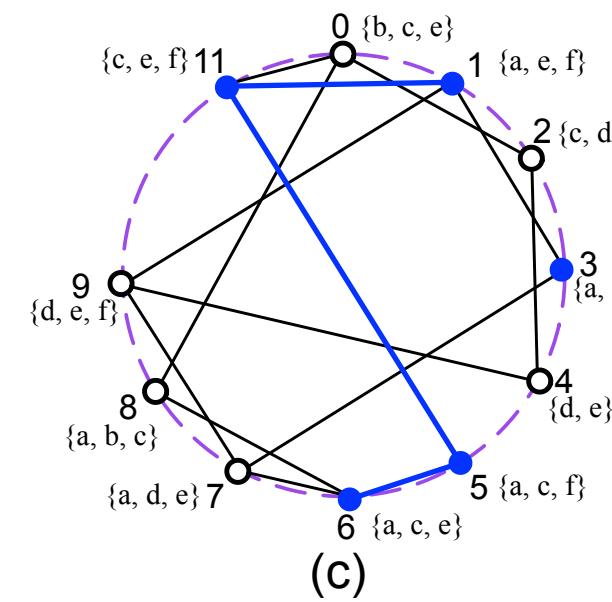
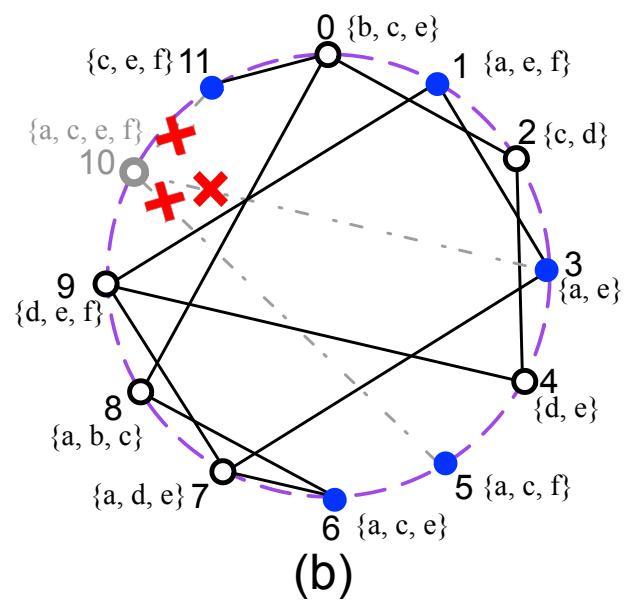
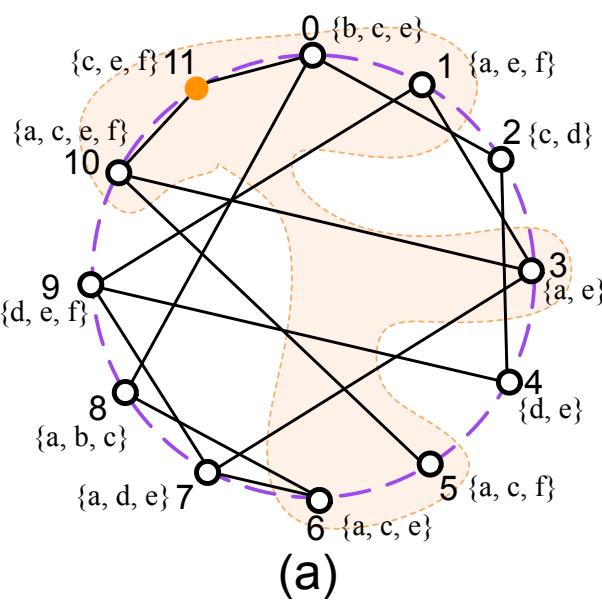


Example: node departure handled

<i>.self</i>	v_{11}
<i>.N</i>	$\{v_0, v_{10}\}$
<i>.D</i>	\emptyset
<i>.succ[1]</i>	v_0
<i>.pred[1]</i>	v_{10}
<i>.pneighb[1]</i>	$\{v_3, v_5, v_{11}\}$
<i>.backup[1]</i>	$\{v_1, v_6\}$

<i>.self</i>	v_{11}
<i>.N</i>	$\{v_0, v_{10}\}$
<i>.D</i>	$\{v_{10}\}$
<i>.succ[1]</i>	v_0
<i>.pred[1]</i>	v_{10}
<i>.pneighb[1]</i>	$\{v_3, v_5, v_{11}\}$
<i>.backup[1]</i>	$\{v_1, v_6\}$

<i>.self</i>	v_{11}
<i>.N</i>	$\{v_0, v_1, v_5\}$
<i>.D</i>	$\{\}$
<i>.succ[1]</i>	v_0
<i>.pred[1]</i>	v_9
<i>.pneighb[1]</i>	$\{v_1, v_4, v_7\}$
<i>.backup[1]</i>	$\{v_5, v_7\}$



Evaluation: experimental setup

- Real-world pub/sub workloads



- Synthetic pub/sub workloads

Expo, Zipf, Unif

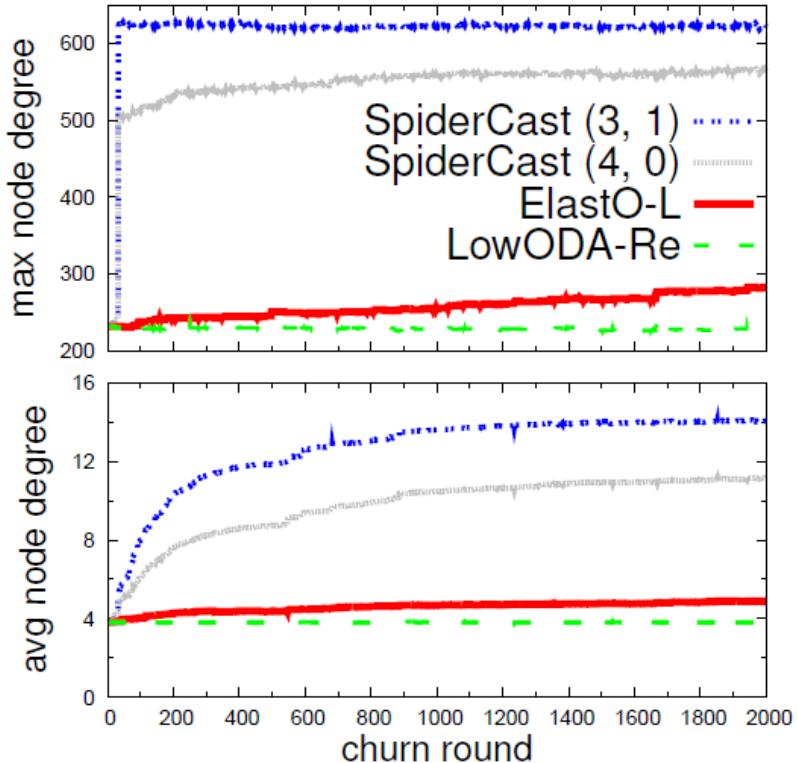
- Churn traces



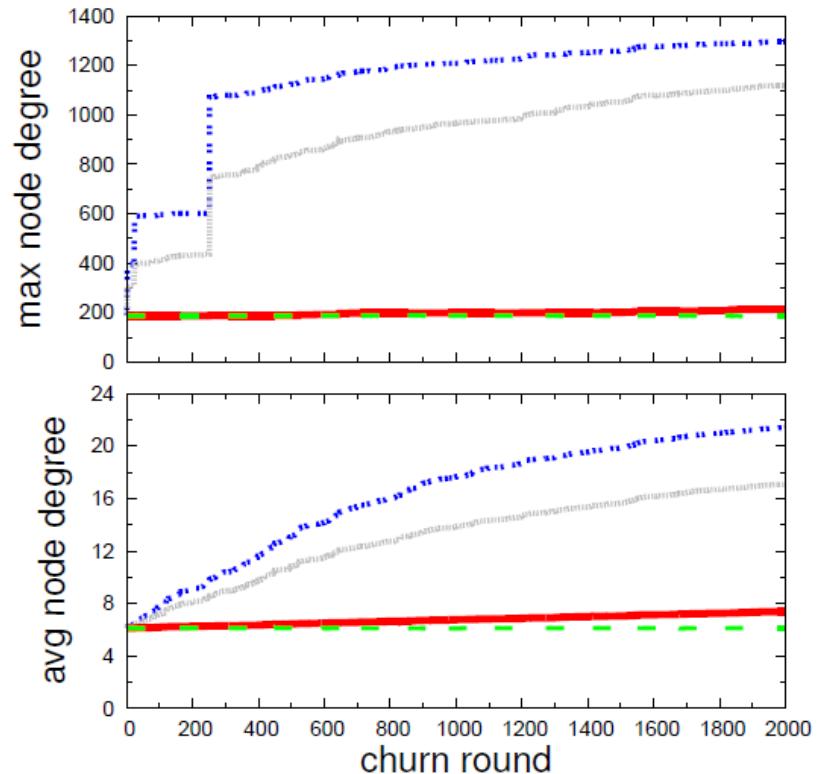
Evaluation: algorithms and protocols

	Our proposed system
ElastO-L	Local view at each node
ElastO-G	Global view at each node
<hr/>	
LowODA	Low Max and Avg Degree Overlay Design Algorithm
LowODA-Inc	Incrementally repair TCO regarding existing links
LowODA-Re	Reconstruct TCO from scratch regardless of existing links
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SpiderCast	A peer-to-peer protocol to build TCO
SpiderCast(Kg,Kr)	Two neighbor selection heuristics: greedy and random

Evaluation: node degrees under churn

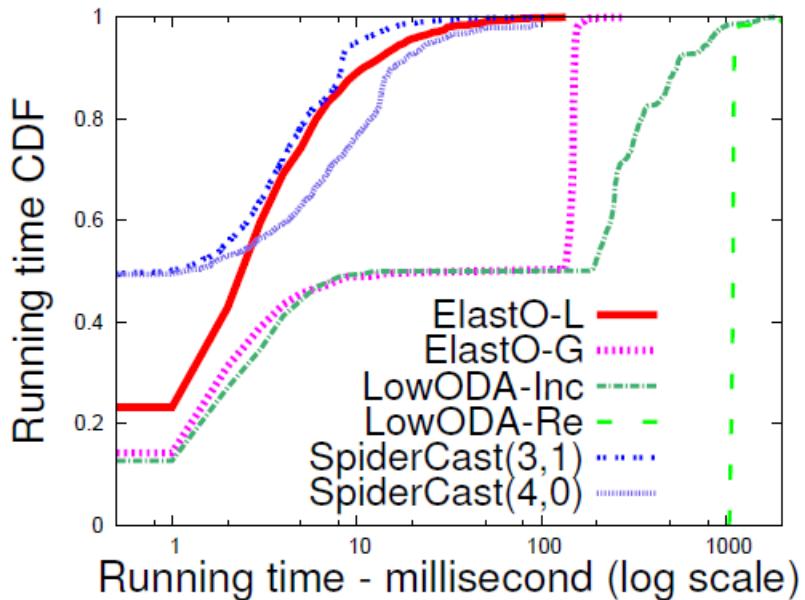


(a) FB 1K

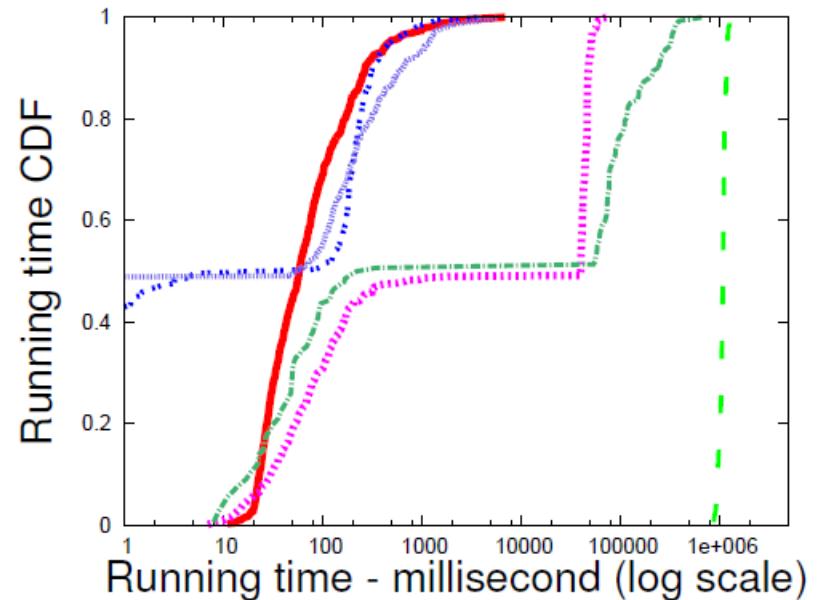


(b) FB 10K

Evaluation: runtime cost



(c) TW 1K



(d) TW 10K

Conclusion

		Knowledge	Churn Handling	Time	Average degree	Maximum degree
Hybrid	ElastO	Local	✓	Fast	$\approx O(\rho \ln V T)$	$\approx O(V /\rho \ln V T)$
Centralized algorithms	Low-ODA	Global	✗	Slow	$O(\rho \ln V T)$	$O(V /\rho \ln V T)$
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	GM ₂				$O(U + \ln V T)$	$\Theta(V)$
Decentralized protocols	SpiderCast, PolderCast, etc.	Local/Global	✓	Fast	Unknown	Unknown

Current research in IBM

Distributed pub/sub for federated messaging and IoT

- Scalability: 100 Million users/topics/subscriptions
- Quality of Service (QoS) for MQTT
 - 0: at most once
 - 1: at least once
 - 2: exactly once
- Topics and subscriptions
 - hierarchy and wildcards
 - e.g., sensors/+/temperature/+
- Overlay topology and routing protocols
- Membership management